

CSC 240

~~DECIDABILITY~~

TURING LANGUAGES

A Turing machine **accepts** a string if it enters an accept state while processing that string.

A Turing machine **rejects** a string if it enters a reject state while processing that string.

A Turing machine **loops** forever on a string if it never enters an accept or reject state while processing that string.

A Turing machine **halts** on a string if it accepts or rejects that string.

TURING LANGUAGES

A language is “**Turing Recognizable**” if a Turing machine recognizes it.

A language is “**Turing Decidable**” if a Turing machine decides it, that is if the strings of that language cause the Turing machine to end up in an **accept** or **reject** state.

Can we build a Turing Machine that *recognizes* A_{TM} ?

Can we build a Turing Machine that *decides* A_{TM} ?

ACCEPTANCE PROBLEMS

We can build a Turing Machine that tells us:

If a DFA will accept a string.

If an NFA will accept a string.

If a Regular Expression will accept a string.

If a CFG describes a string.

If a PDA accepts a string.

ACCEPTANCE PROBLEMS

M - Some Turing Machine

$\langle M, w \rangle$ - Some Turing Machine and its string input encoded as a string

$$A_{\text{TM}} = \{ \langle M, w \rangle \mid M \text{ is a TM that accepts string } w \}$$

Can we build a Turing Machine that *decides* A_{TM} ?

In other words, can we build a Turing Machine that tells us, for an arbitrary TM, M and an arbitrary input string w, if M will accept or reject w?

UNIVERSAL TURING MACHINE

$\langle M, w \rangle$

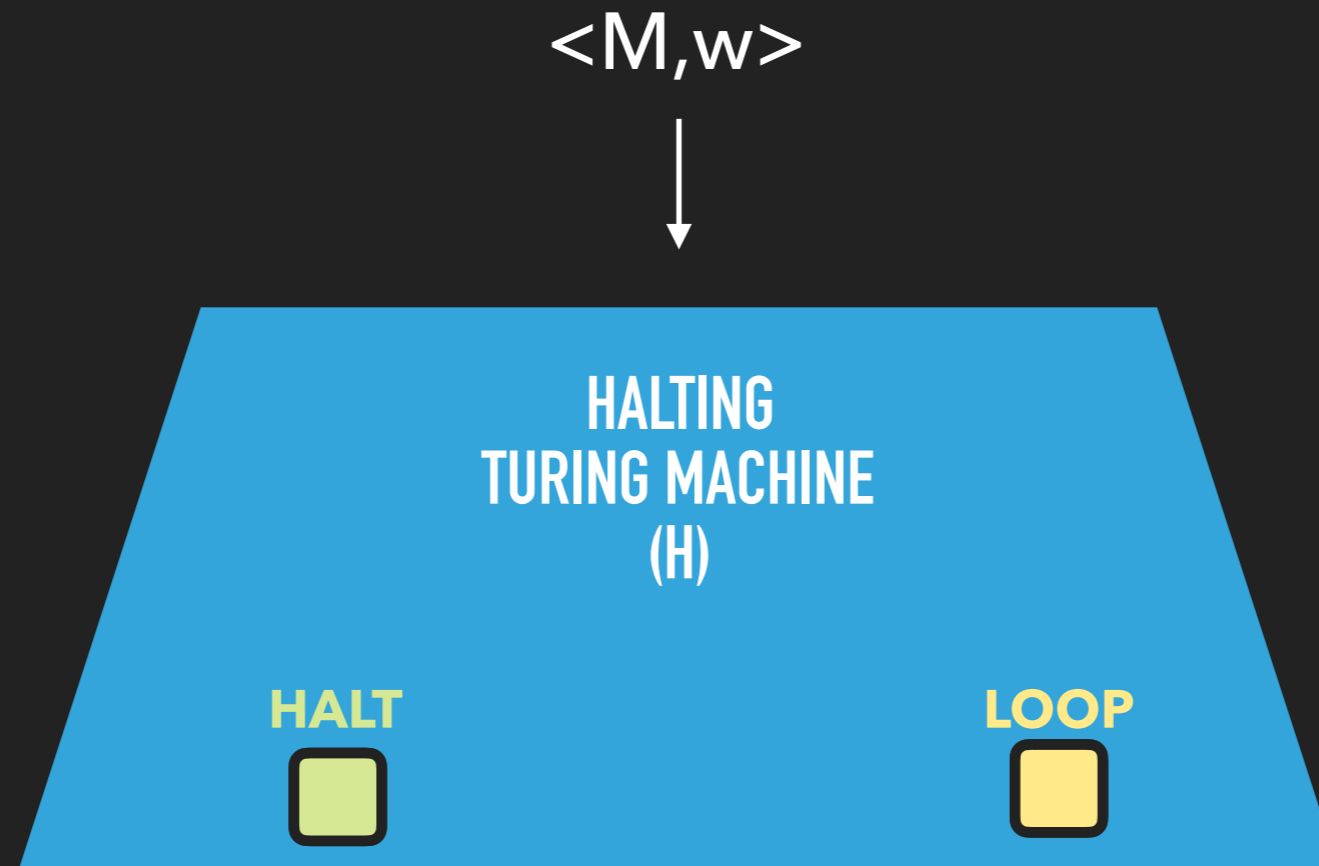


UNIVERSAL
TURING MACHINE
(U)

U is a Universal Turing Machine

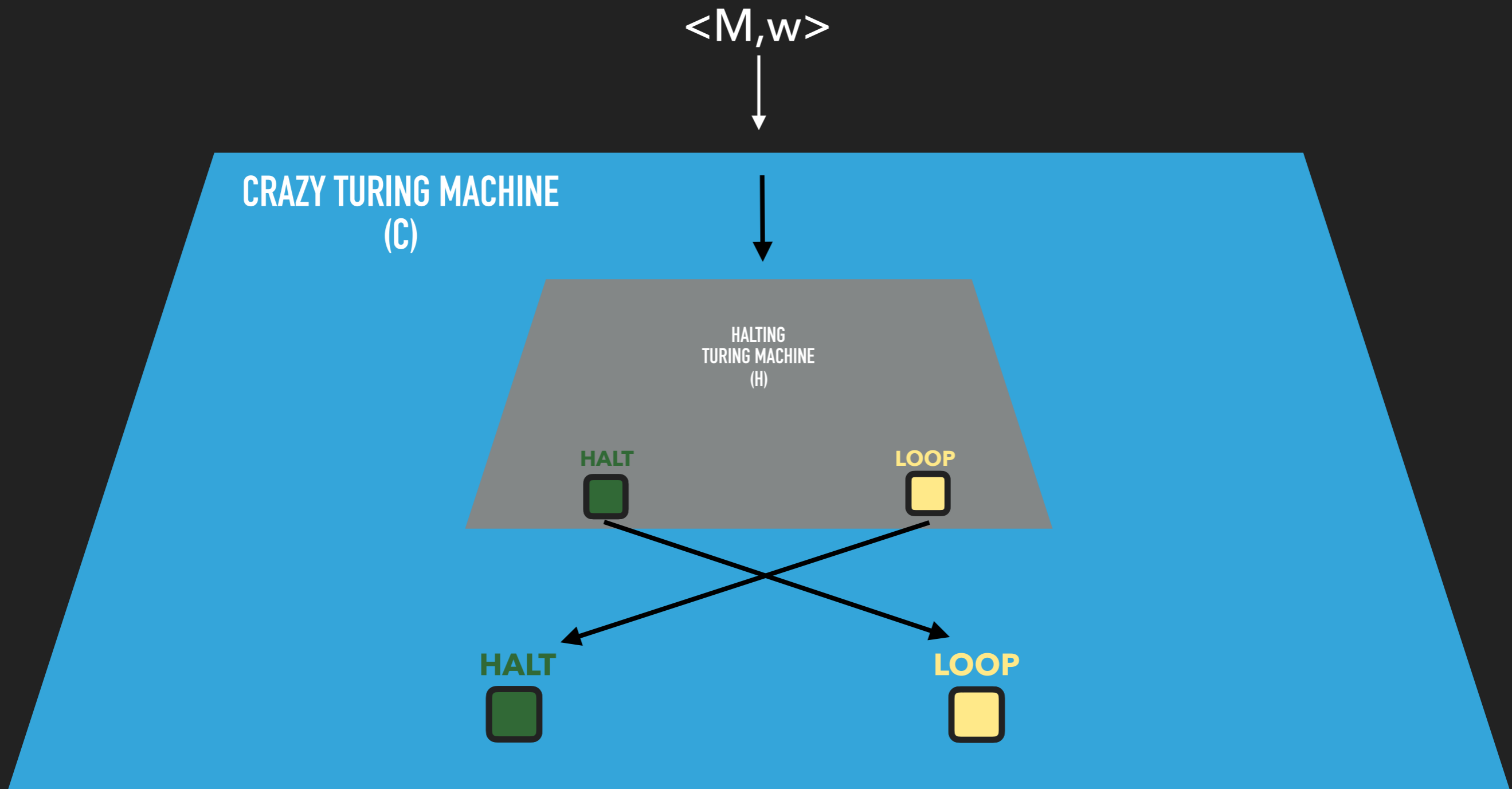
It can accept any Turing Machine and an input string
and run the Turing Machine as it processes w.

THE HALTING MACHINE



If M halts on w , H will halt on $\langle M, w \rangle$
If M loops on w , H will loop on $\langle M, w \rangle$

THE HALTING MACHINE



If M halts on w , C will loop on $\langle M, w \rangle$

If M loops on w , C will halt on $\langle M, w \rangle$

THE HALTING MACHINE

Make C decide if C halts on $\langle C \rangle$

$\langle C \langle C \rangle \rangle$

CRAZY TURING MACHINE
(C)

HALTING
TURING MACHINE
(H)

CONTRADICTION!!!

HALT
■

LOOP
■

If C halts on $\langle C \rangle$, C will loop on $\langle C \rangle$

If C loops on $\langle C \rangle$, C will halt on $\langle C \rangle$

ACCEPTANCE PROBLEMS

$\langle M, w \rangle$ - Some Turing Machine and its string input encoded as a string

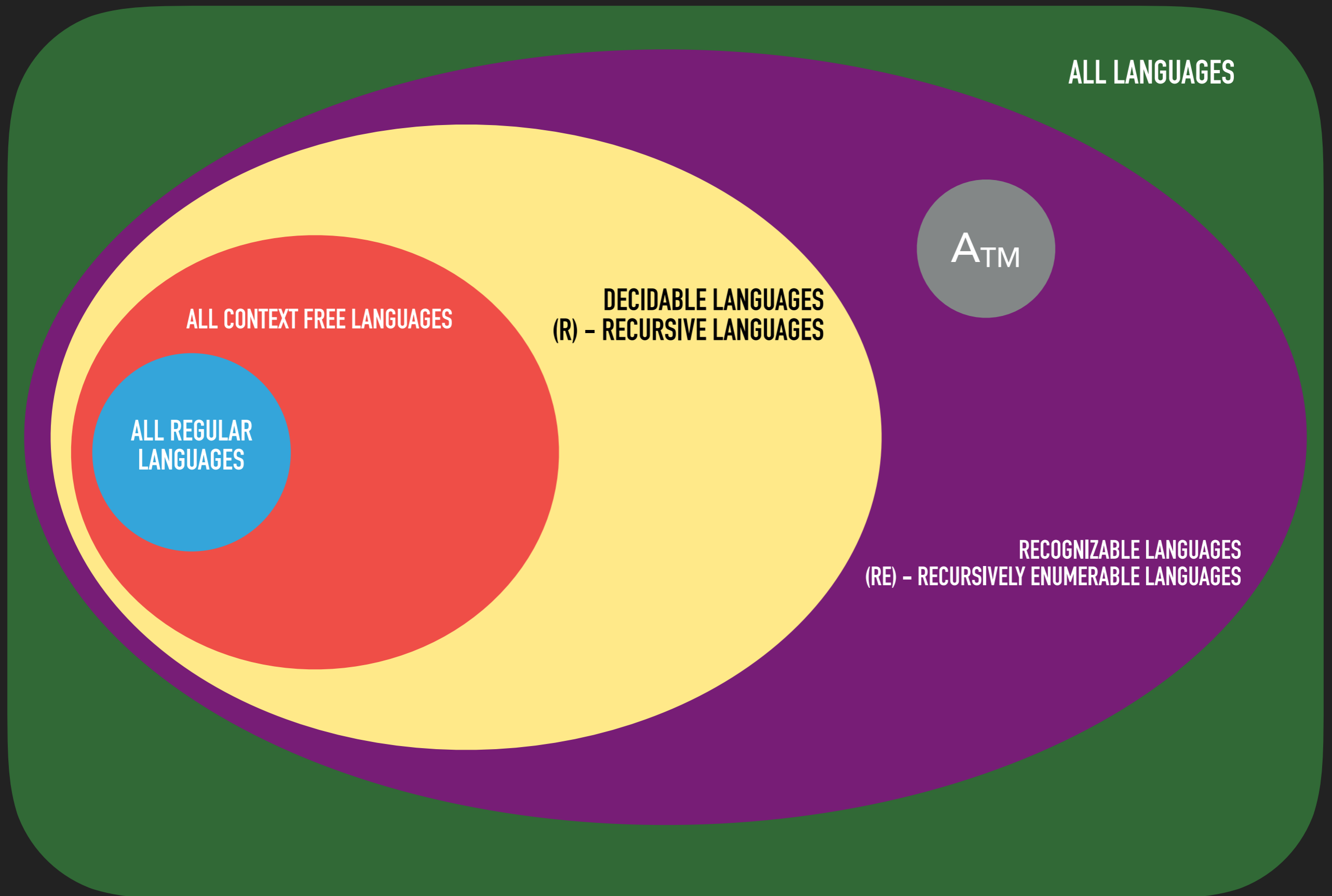
$$A_{TM} = \{ \langle M, w \rangle \mid M \text{ is a TM that accepts string } w \}$$

Can we build a Turing Machine that *decides* A_{TM} ?

In other words, can we build a Turing Machine that tells us, for an arbitrary TM, M and an arbitrary input string w , if M will accept or reject w ?

NO! A_{TM} is recognizable, but not decidable.

TURING DECIDABLE LANGUAGES



ACCEPTANCE PROBLEMS

We can build a Turing Machine that tells us:

If a DFA will accept a string.

If an NFA will accept a string.

If a Regular Expression will accept a string.

If a CFG describes a string.

If a PDA accepts a string.

} Decidable

We can't build a Turing Machine that tells us:

If a TM will halt (accept or reject) on a string. Not Decidable

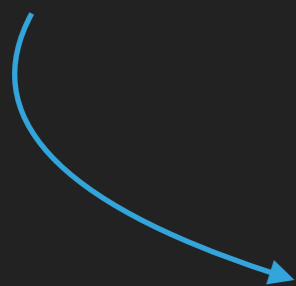
All of these problems (languages) are *recognizable* by at least one TM.

Are there problems (languages) *not recognizable* by any TM?

UNRECOGNIZABLE LANGUAGES

$$\text{TM} = \{M_0, M_1, M_2, M_3, \dots\}$$

M_0 recognizes these strings...



	$\langle M_0 \rangle$	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$...
M_0	Y	N	Y	Y	... , ... }
M_1	N	Y	Y	Y	... , ... }
M_2	Y	Y	N	N	...
M_3	N	N	Y	N	...
...

UNRECOGNIZABLE LANGUAGES

$$\text{TM} = \{M_0, M_1, M_2, M_3, \dots\}$$

	$\langle M_0 \rangle$	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$...
M_0	Y	N	Y	Y	...
M_1	N	Y	Y	Y	...
M_2	Y	Y	N	N	...
M_3	N	N	Y	N	...
...

The "Diagonal Language" L_D

Y	Y	N	N	...
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does this row have a pairing?

UNRECOGNIZABLE LANGUAGES

$$\text{TM} = \{M_0, M_1, M_2, M_3, \dots\}$$

	$\langle M_0 \rangle$	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$...
M_0	Y	N	Y	Y	...
M_1	N	Y	Y	Y	...
M_2	Y	Y	N	N	...
M_3	N	N	Y	N	...
...

Generate the complement of the row by flipping its Y's and N's.

The Complement of the Diagonal Language.

$\overline{L_D}$

N	N	Y	Y	...
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does this row have a pairing?

UNRECOGNIZABLE LANGUAGES

$$\mathbf{TM} = \{M_0, M_1, M_2, M_3, \dots\}$$

N	N	Y	Y	...
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$$M? \leftrightarrow \{ \langle M_3 \rangle, \langle M_4 \rangle, \dots \}$$

This language is not recognizable by any Turing Machine!

There is NO Turing Machine that will recognize this set of strings.

$\overline{L_D}$ is not recognizable by any Turing Machine

$$|\mathbf{TM}| < |\mathcal{L}|$$

UNRECOGNIZABLE LANGUAGES

How many different possible Turing machines are there?

How many different possible Languages are there?

$\overline{L_D}$ is not recognizable by any Turing Machine

$$|\text{TM}| < |L|$$

LANGUAGES

